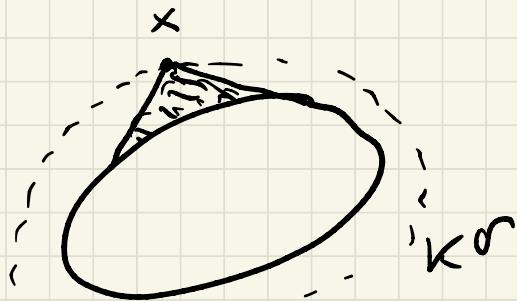


dual concept to floating body

illumination body (w)



$$|[\bar{x}, K] \setminus K| \leq \sigma$$



$$K^\sigma = \{x : |[\bar{x}, K] \setminus K| \leq \sigma\}$$

K^σ is convex

w

$$\lim_{\sigma \rightarrow 0} \frac{|K^\sigma| - |K|}{\sigma^{\frac{2}{n+1}}} = d_n \text{ as } (K)$$

L_p -affine surface

$$-\infty \leq p \leq +\infty$$

$$p \neq -n$$

$$\text{as}_p(K) = \int \frac{\frac{x}{(n+p)}}{\partial K \cdot ((x, N(x)))^{\frac{n(p-1)}{n+p}}} d\mu_K \quad 0 \in \text{int}(K)$$

$$p=1 \rightarrow \text{as}_1(K) = \text{as}(K)$$

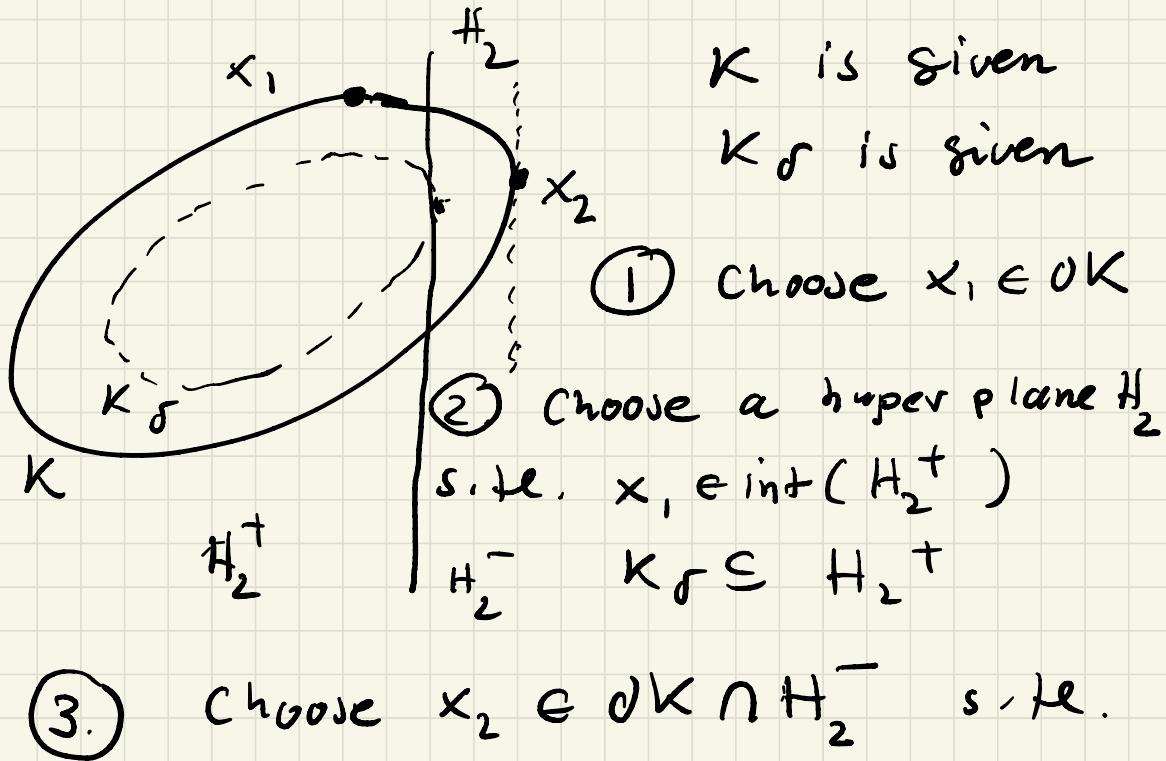
$$p=0 : \text{as}_0(K) = \int_{\partial K} (x, N(x)) \circ d\mu_K \\ = n(K)$$

$$p=\pm\infty : \text{as}_{\pm\infty}(K) = n \underline{|K^\circ|}$$

II. Approximation of convex bodies

by polytopes

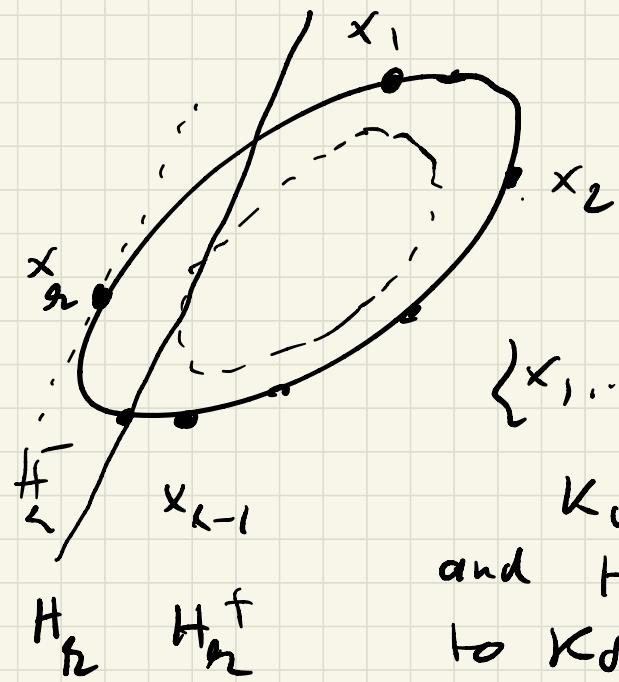
II.1 The floating body algorithm (Schütt)



By choice of H_2^- : $|H_2^- \cap K| \geq \delta$

Suppose we have chosen x_1, \dots, x_{k-1} .

Now choose x_k



we choose a hyperplane H_k

s.t.

$$\{x_1, \dots, x_{k-1}\} \subseteq H_k^+$$

$$K \subseteq H_k^+$$

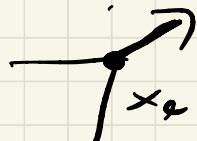
and H_k is a supporting hyperplane

to $K \cap$

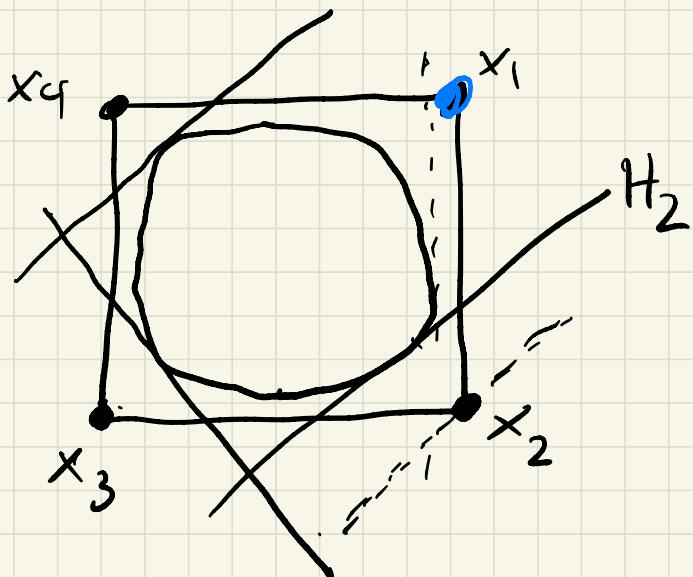
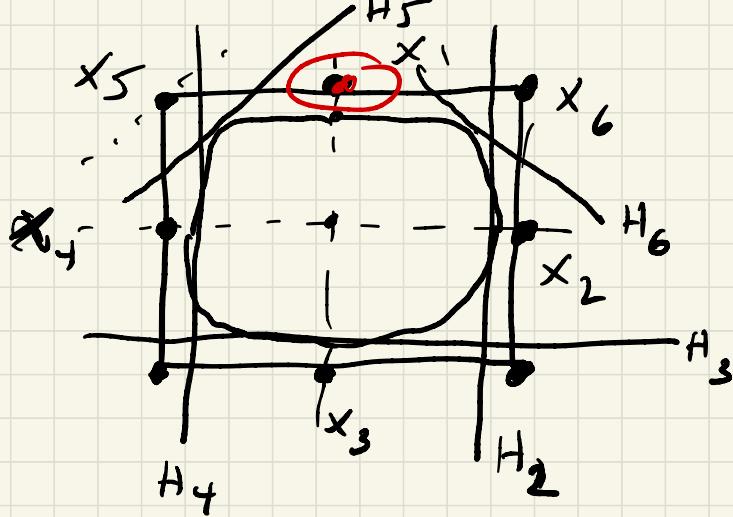
and then we choose a point

$x_k \in K \cap H_k^-$ s.t. $N_k(x_k)$ is

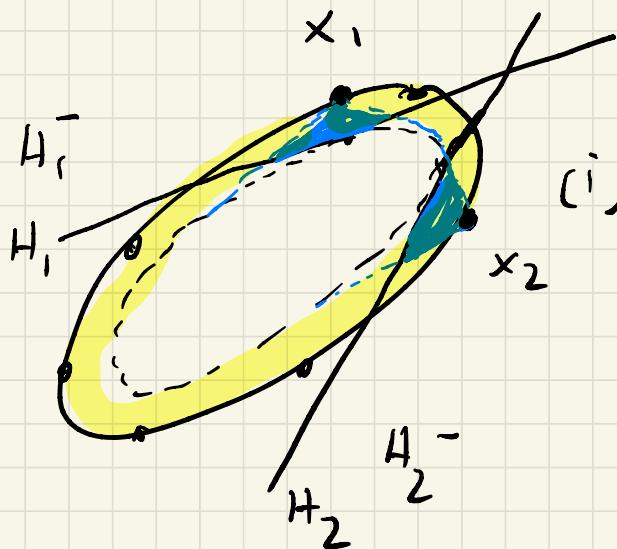
a normal to H_k



Example of square = B_{∞}^2



How many points can we choose?



Note

$$(i) [x_k, K_\delta] \cap H_k^- \cap$$

$$[x_{k+1}, K_\delta] \cap H_{k+1}^- = \emptyset$$

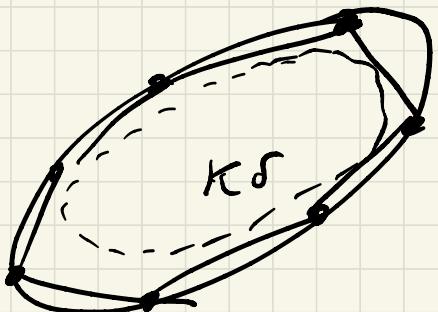
$$(ii) \underbrace{|[x_k, K_\delta] \cap H_k^-|}_{\text{oval}} \approx \delta$$

How many points can we choose?

$$N, \delta = \sum_{k=1}^N \underbrace{|[x_k, K_\delta] \cap H_k^-|}_{\text{line}}$$

$$= \left| \bigcup_{k=1}^N ([x_k, K_\delta] \cap H_k^-) \right|$$

$$\approx |K \setminus K_\delta| \Rightarrow N \approx \frac{|K \setminus K_\delta|}{\delta}$$



$$P_N = [x_1, \dots, x_N]$$

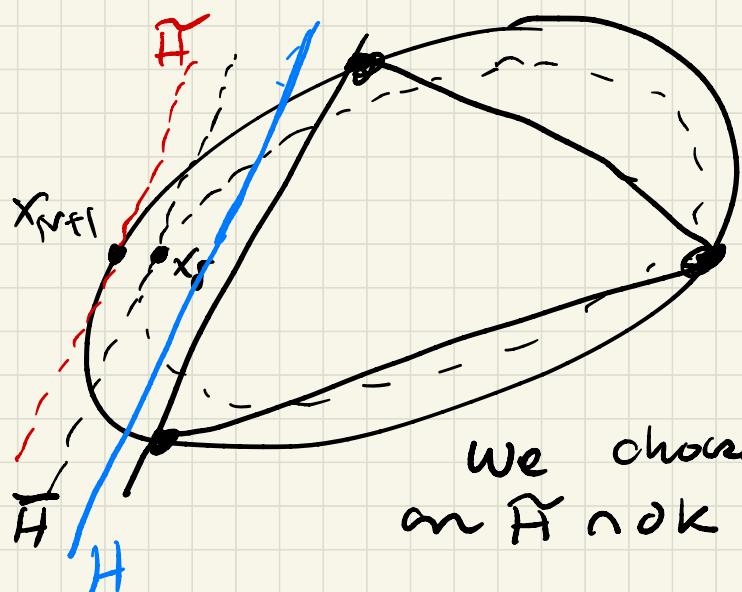
P_N "approximates" K

Note $\boxed{K_\delta \subseteq P_N} \subseteq K$

We show: $K_\delta \subseteq P_N$.

Suppose not. Then there is $x_\delta \in K_\delta$,

$$x_\delta \notin P_N$$



By separation theorem \exists a hyperplane H

i.e.

$$x_\delta \in H^- \text{ and } P_N \subseteq H^+$$

We choose a point x_{N+1} on $H \cap K$

Theorem (Schütt)

For all $\delta \leq \frac{|K|}{4e^4}$ there is a NEN s.t.

$$\frac{N}{\pi} \leq e^{16n} \frac{|K \setminus K_\delta|}{\delta |B_2^n|} \quad ||$$

and there is a polytope P_N with at most N vertices s.t.

$$K_\delta \subseteq P_N \subseteq K$$

(we want to see how good this approximation is: Let K be s.t. $as(K) > 0$)

$$\frac{|K \setminus K_\delta|}{\delta^{\frac{2}{n+1}}} \xrightarrow{\delta \rightarrow 0} \frac{1}{2} \left(\frac{n+1}{|B_2^{n+1}|} \right)^{\frac{2}{n+1}} \underbrace{as(K)}_{\approx C \cdot n \underbrace{as(K)}_{\text{---}}}$$

By the theorem there $N \approx \frac{e^{(K)K_0}}{\delta^{(B_2^n)}}$

and a polytope P_N will have at most N vertices or be. $K_f \subseteq P_N \subseteq K \Rightarrow$

$$\frac{|K \setminus P_N|}{N^{-\frac{2}{n-1}}} \leq \frac{|K \setminus K_f|}{N^{-\frac{2}{n-1}}} \approx \frac{|K \setminus K_f|^{\frac{n+1}{n-1}}}{\delta^{\frac{2}{n-1}} (B_2^n)^{\frac{2}{n-1}}} e^{\frac{2n}{n-1}}$$

$$= \left(\frac{|K \setminus K_f|}{\delta^{\frac{2}{n+1}}} \right)^{\frac{n+1}{n-1}} \cdot c \cdot n$$

$$\leq \underline{c} n^2 \omega(K)^{\frac{n+1}{n-1}}$$

How good is this result?

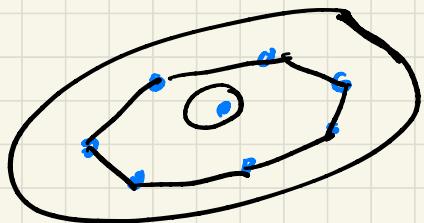
II. 1 Random polytopes

A random polytope is the convex hull of N randomly iid chosen points in \mathbb{R}^n with respect to a probability measure

$$P_N = [x_1, \dots, x_N]$$

typical settings

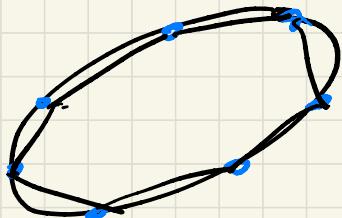
- (i) N points are chosen at random w.r.t. $\frac{m}{|K|}$ in a convex body K



→ not all points chosen become vertices of P_N

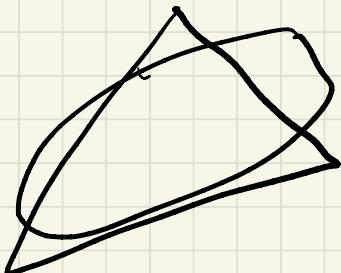
we have wasted points

(ii) Choose the points on ∂K
instead



The setting (i), (ii) produces approximating polygons inside K

(iii) arbitrary positioned polygon



Case (ii)

We will choose N points on the boundary of a convex body K

$$\text{with respect to } P_f = f \cdot d\mu$$

where f is a continuous strictly positive function s.t. $\int_K f d\mu_K = 1$

The expected volume of such a random polygon is

$$\underline{\mathbb{E}_N} (\partial K, P_f) = \overbrace{\int \dots \int}^{\partial K} |[x_1 \dots x_N]|$$

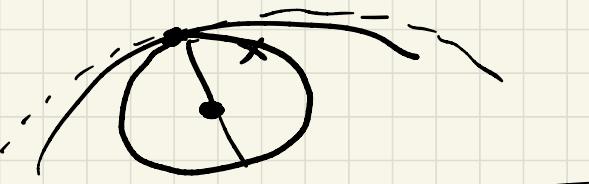
$$dP_f(x_1) \dots dP_f(x_N)$$

Theorem (Schütt, W.; Reiber; K. Knutson; D. Müller)

Let K be a convex body in \mathbb{R}^n s.t.

$\exists 0 < r \leq R < \infty$ s.t. $\forall x \in \partial K$

$$B_2^n(x - rN(x), r) \subseteq K \subseteq B_2^n(x - RN(x), R)$$



Then



$$\lim_{N \rightarrow \infty} \frac{|K| - E_N(\partial K, P_f)}{N^{-\frac{2}{n-1}}} = d_n \int_{\partial K} \frac{x \, d\mu_K}{f^{\frac{2}{n-1}}}$$

Remarks

- we know d_n explicitly

- When the points are chosen in K

$$|K| - E_N(K, \frac{m}{|K|}) \xrightarrow{N \rightarrow \infty} \beta_n \text{as}(K) \left(\frac{|K|}{N} \right)^{\frac{2}{n-1}}$$

Q: $\underset{f \in \mathcal{F}}{\text{Which } f \text{ is the RHS minimal}}$

A: This happens w/ $f_\alpha(x) = \frac{x^{\frac{1}{n+1}}}{\alpha s(k)}$



$$\lim_{N \rightarrow \infty} \frac{|K| - \mathbb{E}_N(\delta K, \mathbb{P}_{f_\alpha})}{N^{-\frac{2}{n-1}}} = d_n$$

$$\int \frac{x^{\frac{1}{n+1}} d\mu_K}{(K^{\frac{1}{n+1}})^{\frac{2}{n-1}}} \underset{as(K)^{\frac{2}{n-1}}}{\approx}$$

$$= d_n \underset{as(K)^{\frac{n+1}{n-1}}}{\approx}$$

Floating body algorithm:

$$\frac{|K \setminus P_N|}{N^{-\frac{2}{n-1}}} \leq c \cdot n^2 \underset{as(K)^{\frac{n+1}{n-1}}}{\approx}$$

One can show: $d_n \approx n$

Random reset

$$\frac{|K| - \mathbb{E}_N(\partial K, P_{f_\alpha})}{N^{-\frac{2}{n-1}}} \underset{N \rightarrow \infty}{\approx} \frac{d_n}{as(K)^{\frac{n+1}{n-1}}}$$

Compare this to the best approximating reset

Theorem (McClure, Vitale $n=2$
Gruber $n \geq 3$)

K is C^2_+

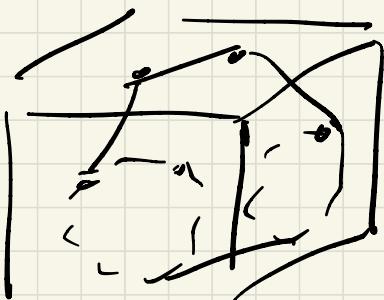
$$\frac{|K| - |P_N^{best}|}{N^{-\frac{2}{n-1}}} \underset{N \rightarrow \infty}{\approx} \frac{\alpha_n}{as(K)^{\frac{n+1}{n-1}}}$$

$$\alpha_n \leq d_n \leq \alpha_n \left(1 + \frac{\log n}{n}\right)$$

Mankiewicz, Schiltz $\alpha_n \approx n$

Reiter, Schütt, W

What happens when we choose
points at random on the
boundary of polygon



Thank you!