

Decision tree:

$$T = (E, \{\varphi_t\}_{t \in E})$$

For $f: \{0, 1\}^E \rightarrow \mathbb{R}$, define

$$\tau(w) = \tau_{f, T}(w)$$

$$:= \min \{t \geq 1 : \forall w' \in \{0, 1\}^E$$

$$\text{if } w'_{[e_t]} = w_{[e_t]} \text{ then } f(w') = f(w)\}$$

this means that by time t
the value of f is determined.

For convenience, we will be
defining τ also beyond $\tau_T(w)$

Lecture 11

Lemma 1

$$\frac{d}{dp} \varphi_{(2n), p, q}^1 (0 \in \partial A_n)$$

$$= \frac{1}{p(1-p)} \sum_{e \in E} \varphi_{(2n), p, q}^1 (1 \in \partial A_n, w_e)$$

Proof:

Below we fix Λ, p, q and omit them in the writing. For a rand. var. X , we define:

$$Z(X) := \sum_{\omega \in \{0,1\}^{\mathbb{N}}} X(\omega) \cdot \underbrace{p^{o(\omega)}}_{\sum_e \omega_e} \underbrace{(1-p)^{c(\omega)}}_{\sum_e (1-\omega_e)} q^{k(\omega)}$$

In particular: $Z(1) = \underbrace{Z}_{\text{partition func.}}$

Then,

$$\frac{d}{dp} Z(X) = \sum_{\omega \in \{0,1\}^{\mathbb{N}}} X(\omega)$$

$$= \left[\sum_e \omega_e \cdot \frac{1}{p} \cdot p^{o(\omega)} (1-p)^{c(\omega)} q^{k(\omega)} \right.$$

$$\left. - \sum_e (1-\omega_e) \cdot \frac{1}{1-p} \cdot p^{o(\omega)} (1-p)^{c(\omega)} q^{k(\omega)} \right]$$

$$= Z \cdot \sum_e \mathbb{E} \left(X \cdot \left[\frac{\omega_e}{p} - \frac{1-\omega_e}{1-p} \right] \right)$$

$$= Z \cdot \frac{1}{p(1-p)} \cdot \sum_e \mathbb{E} \left(X \cdot \underbrace{[\omega_e - \omega_e p - p + p \omega_e]}_{\omega_e - p} \right)$$

Note that $\mathbb{E}(X) = \frac{Z(X)}{Z}$

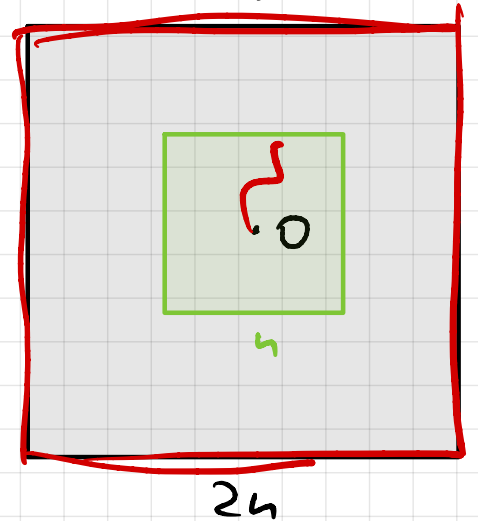
then,

$$\begin{aligned} \frac{d}{dp} E(X) &= \frac{d}{dp} \left(\frac{z(X)}{z(1)} \right) = \frac{z(X)'}{z} - \frac{z(1)'}{z^2} \cdot E(X) \\ &= \frac{1}{p(1-p)} \underbrace{\sum_e \left[E(X \cdot (\omega_e - p)) - E(\omega_e - p) E(X) \right]}_{\text{Cov}(X, \omega_e)} \end{aligned}$$

Take $X = \mathbb{1}_{0 \in \partial \Lambda_n}$.

$$\Theta_n(p) = \varphi_{\Lambda_{2n}, p, q}^1(0 \in \partial \Lambda_n)$$

$$S_n(p) = \sum_{k=0}^{n-1} \Theta_k(p)$$



Lemma 2

$$\underbrace{\Theta_n(1 - \theta_n)}_{\leq} C_{\Lambda_{2n}, p, q} \cdot \frac{8 S_n}{n}$$

$$\text{Var}(\mathbb{1}_{0 \in \partial \Lambda_n}) \cdot \sum_e \text{Cov}(\mathbb{1}_{0 \in \partial \Lambda_n}, \omega_e),$$

$$C_{G, p, q} := \left[\min_{e \in E_{G, p, q}} (\text{Var}(\omega_e)) \right]^{-1}$$

Thm (OSSS inequality)

Let $q \geq 1$, $p \in [0, 1]$, G - finite graph.
Let $f: \{0, 1\}^E \rightarrow [0, 1]$ be increasing.
Fix a decision tree T for f .
Then,

$$\text{Var}(f) \leq C_{G, p, q} \sum_{e \in E} \text{Cov}(f, w_e) \cdot \delta_e(f, T)$$

where

$$\delta_e(f, T) := \varphi_{G, p, q}^{\frac{1}{q}} \left(\underbrace{\exists t \leq \tau(\omega) : e_t = e}_{\text{we reveal } e \text{ to know } f(\omega)} \right)$$

revelment of e for T .

$C_{G, p, q}$ - as in Lemma 2.

Originally found by
O'Donnell, Saks, Schramm, Servedio '05
for Bernoulli percolation.

In 2017, it was extended
to the FK percolation:
Deamuel-Capria, Garofali, Tassion

We'll prove it next time.

To prove Lemma 2, we need to find a tree with a small revealment - on every edge!

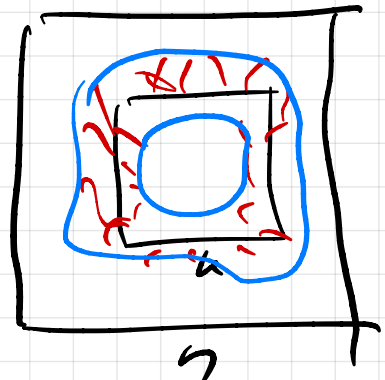
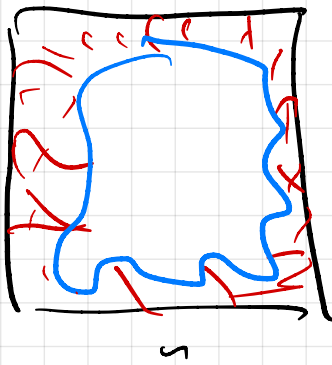
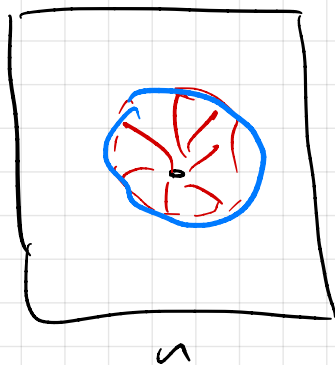
Wrong idea 1:

Explore all edges - revealment for every edge will be high.

Wrong idea 2:

Explore edges from the origin, more efficient - we explore less edges.

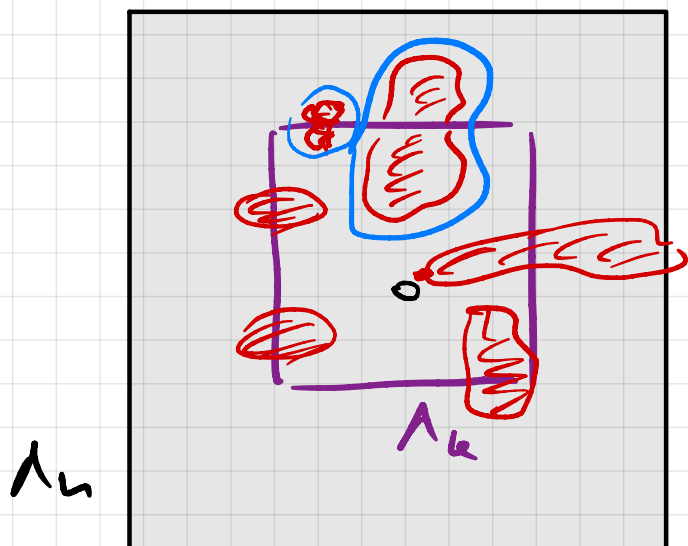
Problem - next to the origin we'll explore everything. In the rest - very good.



Idea: average over all scales!

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Proof:



Decision tree T_k :
explore all clusters crossing

$\partial \Lambda_k$

If one of the contains o

and crosses $\partial \Lambda_n$,

then $o \in \partial \Lambda_n$.

Otherwise - there is a blocking circuit and $o \notin \partial \Lambda_n$.

It's enough to explore only edges $e = uv$, s.t.

$u \in \partial \Lambda_k$ or $v \in \partial \Lambda_k$.

$$d_{uv}(\mathbb{1}_{o \in \Lambda_n}, T_k)$$

$$\leq \varphi_{\Lambda_{2r}, p, q}^{\perp}(u \in \partial \Lambda_k) + \varphi_{\Lambda_{2r}, p, q}^{\perp}(v \in \partial \Lambda_k)$$

By the OSSS ineq. applied for $f = \mathbb{1}_{o \in \Lambda_n}$ and T_k :

$$\Theta_n(1 - \Theta_n) \leq C \cdot \sum_{e \in E} C_{uv}(\mathbb{1}_{o \in \Lambda_n}, \omega_e) \cdot \left(\varphi_{\Lambda_{2r}, p, q}^{\perp}(u \in \partial \Lambda_k) + \varphi_{\Lambda_{2r}, p, q}^{\perp}(v \in \partial \Lambda_k) \right)$$

Average over all $k=1, 2, \dots, n$.
 On the RHS, in the revelation
 we'll get:

$$\frac{1}{n} \sum_{k=1}^n \varphi_{2n}^{\Delta} (u \leftrightarrow \partial \Delta_k)$$

$$\approx \varphi_{2n}^{\Delta} (u \leftrightarrow \partial \Delta_{\frac{(u)}{\|u\|}})$$

$$\approx \frac{1}{n} \sum_{m=0}^{n-1} \varphi_{2n}^{\Delta} (u \leftrightarrow \partial \Delta_m(u))$$

$$\approx \frac{1}{n} \sum_{m=0}^{\lfloor \frac{n}{2} \rfloor} \varphi_{2n}^{\Delta} (u \leftrightarrow \partial \Delta_m(u))$$

$$\approx \varphi_{2n}^{\Delta} (u \leftrightarrow \partial \Delta_m(u))$$

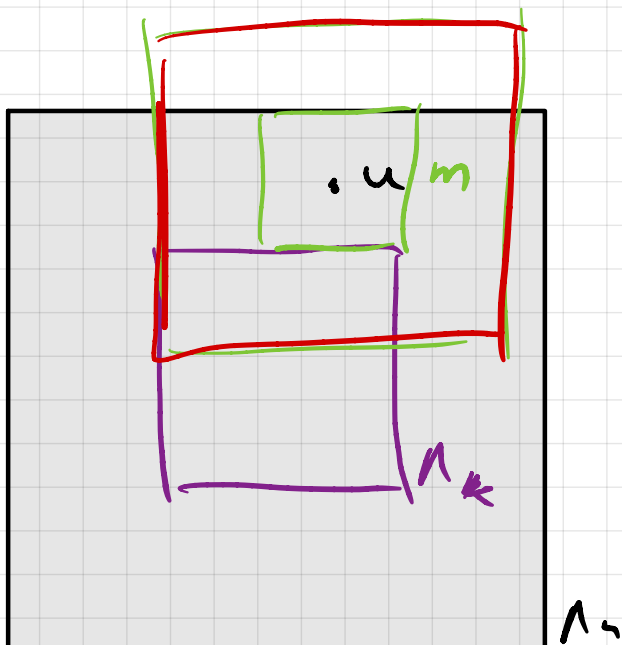
$$\approx \varphi_{2n}^{\Delta} (0 \leftrightarrow \partial \Delta_m)$$

$$\approx \partial \Delta_m$$

In conclusion:

$$\frac{1}{n} \sum_{k=1}^n \varphi_{2n}^{\Delta} (u \leftrightarrow \partial \Delta_k)$$

$$\approx \frac{1}{n} \sum_{k=1}^n \partial \Delta_k$$



Δ_{2n}

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Average revelation:

$$\frac{1}{n} \sum_{k=1}^n (\varphi_{2n}^1(u \leftrightarrow \partial A_k) + \varphi_{2n}^2(v \leftrightarrow \partial A_k))$$

By $\frac{1}{n} \cdot S_n$ and the inequality follows □

Proof (Sharpness given OSSS):

Insert the equality of Lem 1 into Lem 2:

$$\Theta_n'(p) \geq \Theta_n \cdot \frac{1}{S_n}$$

$$\cdot (1 - \Theta_n) \frac{1}{p(1-p)} \cdot \min_{e \in E} (\text{Var}(w_e))$$

$$\geq c(p_0, p_1) > 0$$

whenever $p \in (p_0, p_1)$.

$$0 < p_0 \leq p_1 < 1$$

Thus, we get

$$\Theta_n' \geq c \cdot \frac{1}{S_n} \cdot \Theta_n$$

Note first that

$$\lim_{n \rightarrow \infty} \Theta_n \stackrel{?}{=} \lim_{n \rightarrow \infty} \varphi_{p,q}^{\Delta} (0 \leftrightarrow \partial \Lambda_n) \\ = \varphi_{p,q}^{\Delta} (0 \leftrightarrow \infty).$$

Also, for every $k \geq \Delta$,

$$\lim_{n \rightarrow \infty} \Theta_n \stackrel{?}{=} \lim_{n \rightarrow \infty} \varphi_{n, p, q}^{\Delta} (0 \leftrightarrow \partial \Lambda_k) \\ = \underbrace{\varphi_{p,q}^{\Delta} (0 \leftrightarrow \partial \Lambda_k)}_{\substack{\geq \\ \varphi_{p,q}^{\Delta} (0 \leftrightarrow \infty)}}$$

Hence,

$$\lim_{n \rightarrow \infty} \Theta_n = \varphi_{p,q}^{\Delta} (0 \leftrightarrow \infty)$$

The theorem follows from:

Lemma 3

Let $f_n : [0, x_0] \rightarrow [0, \ell]$ be increasing, differentiable and satisfy:

- $f_n \rightarrow f$

- $f_n' \geq \frac{1}{M_n} \cdot f_n$, where $\sum_{k=0}^{n-1} f_k$

Then, there exists $x_1 \in (0, x_0]$,
 s.t. (i) for $x < x_1$, there exists
 $c_x > 0$, s.t. for any n
 large enough:

$$f_n(x) \leq \exp(-c_x n)$$

(ii) for $x > x_1$,

$$f(x) \geq x - x_1.$$

This lemma can be applied
 for $f_n := \mathbb{E}[\Theta_n]$ and implies:

• $p > p_c$:

$$\varphi_{p, q}^{\Delta} (0 \rightarrow \infty) \geq c(p - p_c)$$

• $p < p_c$:

$$\Theta_n \leq \exp(-c_p n)$$

Since

$$\Theta_n = \varphi_{\Lambda_{2n}, p, q}^{\Delta} (0 \rightarrow \partial \Lambda_n)$$

$$\geq \varphi_{\Lambda_{2n}, p, q}^{\Delta} (0 \rightarrow \partial \Lambda_{2n})$$

This proves the theorem.